# Algorithms on Graphs

Umang Bhaskar & Juhi Chaudhary

Day 2, session 2: TSP



STCS Vigyan Vidushi 2024



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Better algorithms? Can do  $O(2^n)$  by a dynamic program (Held-Karp algorithm)

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Instead of an <u>optimal</u> solution, we look for a <u>near-optimal</u> solution.

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A 2-approximate algorithm gives a cycle of length at most  $2 \times OPT(I)$ .

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We will give a 1.5-approximate algorithm for TSP.

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a c f h

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<u>Claim</u>: length of MST  $\leq$ length of shortest cycle that visits all pts.

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By doubling edges of MST, can get a <u>walk</u> of length  $2 \times MST \leq 2 \times OPT$ 

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A 2-approximation algorithm:

- 1. Find MST (length  $\leq$  OPT)
- 2. Double edges, find Eulerian walk (length  $\leq 2 \times OPT$ )
- 3. Short-circuit edges to get cycle of length  $\leq$  2 x OPT

(Remainder on board)